

Principles of Wireless Sensor Networks

https://www.kth.se/social/course/EL2745/

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Course content

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Where we are

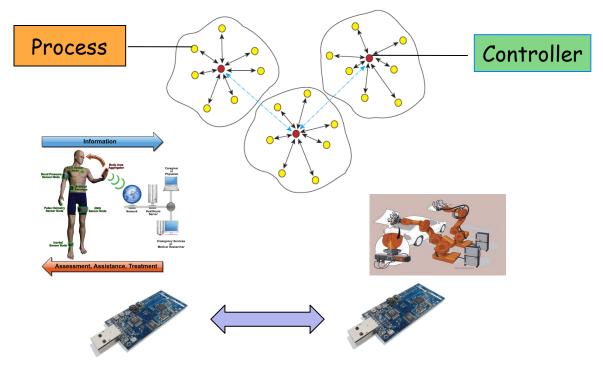
Application

Presentation Session

Transport

Routing

MAC Phy



- How messages are successfully transmitted and received over the wireless channel?
- Aim: modeling the probability to successfully receive messages as function of the radio power, modulations, coding, and channel attenuations normally experienced in WSNs

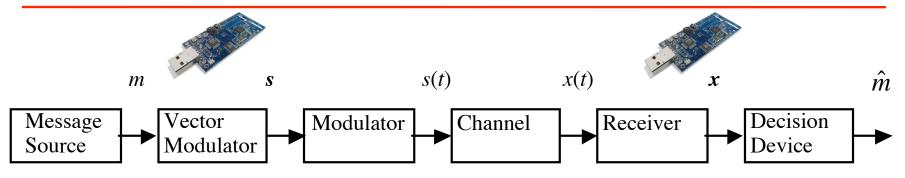
Today's learning goals

- How bits of messages are transmitted over a channel?
- What is the probability to successfully receive messages over AWGN channels?
- What is the probability to successfully receive messages over fading channels?

- Basic of modulation theory
- Probability of error over AWGN channels
 - BPSK
 - Amplitude Modulation
- Probability of error over fading + AWGN channels



Digital modulations



$$s(t) = \sum_{k=0}^{\infty} a_k(t)g(t - kT_s)$$

modulated signal transmitted by the sensor's antenna

$$G(f) \triangleq \int_0^{T_s} a_0 g(t) e^{-2\pi f t} dt$$

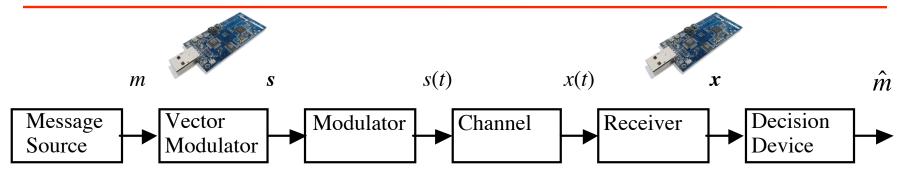
spectrum of the signal over a symbol

$$\Phi_s(f) \triangleq \sigma_a^2 \frac{|G(f)|^2}{T_s}$$

power spectral density of the signal



Example: Binary Phase Shift Keying (BPSK) modulation



$$s(t) = \sum_{k=0}^{\infty} a_k(t)g(t - kT_s)$$

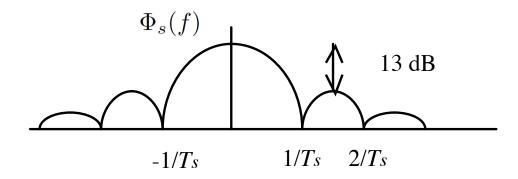
$$a_k(t) = \begin{cases} \cos(2\pi f_c t) & \text{if bit 0 at symbol time k,} \\ \cos(2\pi f_c t + \pi) & \text{if bit 1 at symbol time k.} \end{cases}$$

Modulation of the bits

$$g(t) = \sqrt{\frac{E}{T_s}}, 0 \le t \le T_s$$

$$P_t = \frac{E}{T_s}$$

BPSK spectral density



$$G(f) = -\sqrt{\frac{E}{T_s}} \frac{e^{j\pi f T_s} - e^{-j\pi f T_s}}{j2\pi f} = -\sqrt{\frac{E}{T_s}} T_s \frac{\sin(\pi f T_s)}{\pi f T_s} = -\sqrt{\frac{E}{T_s}} T_s \text{sinc}(f T_s)$$

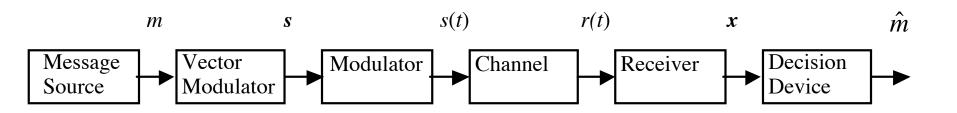
$$\Phi_s(f) \triangleq \frac{E}{T_s} \operatorname{sinc}^2(fT_s)$$



Now, we would like to compute the probability that a bit is received successfully when it is transmitted by a modulation over a AWGN channel

- Basic of modulation theory
- Probability of error over AWGN channels
 - BPSK
 - Amplitude Modulation
- Probability of error over fading + AWGN channels
- Basic of coding theory

BPSK probability of error in AWGN wireless channels



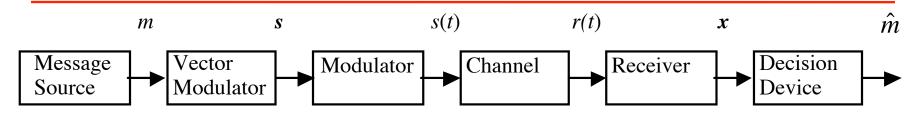
Assume that the transmitted signal is received together with an Additive White Gaussian Noise (AWGN)

$$r(t) = s(t) + n_0(t)$$

$$n_0(t) \in N\left(0, \sigma^2 = \frac{N_0}{2T_s}\right)$$

The demodulator in the receiver produces a signal

BPSK detection in AWGN wireless channels



The demodulator in the receiver block produces a signal

$$r'(t) = s'(t) + n'_0(t)$$

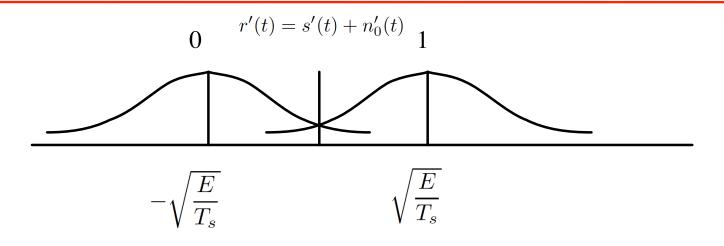
$$s'(t) = \begin{cases} \sqrt{\frac{E}{T_s}} & \text{if bit 0 was transmitted,} \\ -\sqrt{\frac{E}{T_s}} & \text{if bit 1 was transmitted.} \end{cases}$$

$$n_0'(t) \in N\left(0, \sigma^2 = \frac{N_0}{2T_s}\right)$$

- If $r'(t) \ge 0$ the detector decides for bit 1
- If r'(t) < 0 the detector decides for bit 0
- Given the AWGN, what is the error in this detection?



BPSK probability of error

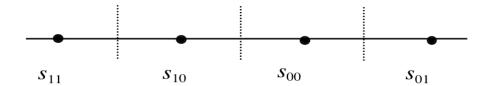


$$P_{e,0|1} = \int_{-\infty}^{0} \frac{1}{\sqrt{2\pi\sigma^2}} e^{-\frac{\left(x - \sqrt{\frac{E}{T_s}}\right)^2}{2\sigma^2}} dx = Q\left(\sqrt{\frac{2E}{N_0}}\right)$$

$$P_{e,1|0} = \int_0^{-\infty} \frac{1}{\sqrt{2\pi\sigma^2}} e^{-\frac{\left(x + \sqrt{\frac{E}{T_s}}\right)^2}{2\sigma^2}} dx = Q\left(\sqrt{\frac{2E}{N_0}}\right)$$

$$P_e = \frac{1}{2} P_{e,0|1} + \frac{1}{2} P_{e,1|0} = Q\left(\sqrt{\frac{2E}{N_0}}\right)$$
 SNR = $\frac{E}{N_0}$

Amplitude modulation, AM



- BPSK is a simple example of binary Amplitude Modulation
- By adding more amplitude values, more general modulations are possible
- Modulations are characterized by constellation points
- Every point (symbol) is associated to a signal with specific amplitude
- Every signal is then associated to a code-word of bits

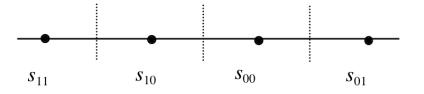
Probability of error for AM

- The distance between the constellation points determines the probability that a symbol is detected with error
- It is possible to show that

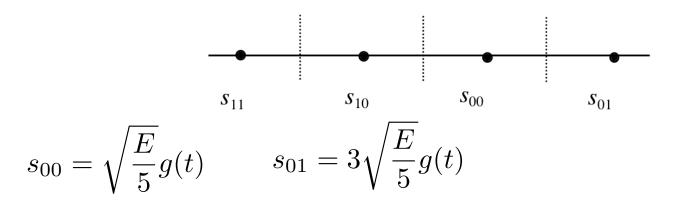
$$P_e \simeq N_{d_{\min}} Q \left(\frac{d_{\min}}{\sqrt{2N_0}} \right)$$

 d_{\min} minimum distance among points

 $N_{d_{\min}}$ average number of neighbors at minimum distance



Exercise: 4-AM (or 4-PAM)



$$s_{11} = -3\sqrt{\frac{E}{5}}g(t)$$
 $s_{10} = -\sqrt{\frac{E}{5}}g(t)$

$$g(t) = \sqrt{\frac{1}{T_s}} \qquad 0 \le t \le T_s$$

What is the probability that a bit is erroneously received?

Solution

$$d_{\min}^2 = 4\frac{E}{5} \qquad N_{\min} = 1.5$$

The probability of symbol error is

$$P_e \simeq 1.5Q \left(\sqrt{\frac{2E}{5N_0}}\right) = 1.5Q \left(\sqrt{\frac{4E_b}{5N_0}}\right)$$

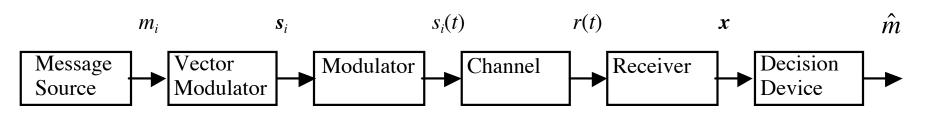
 Since per every symbol there are two bits, the probability of a bit in error is

$$P_b \simeq 0.75Q \left(\sqrt{\frac{4E_b}{5N_0}} \right)$$

$$E_b = \frac{E}{\log_2 M} = \frac{E}{2}$$
 Energy per bit

- Let's consider a more general amplitude modulation that is used in
 - TskyMotes
 - IEEE 802.15.4 standard (which we study next lecture)
- 4-Quadrature Amplitude Modulation





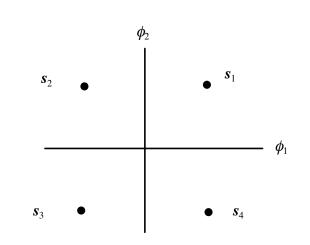
$$s_i(t) = \frac{2E}{T_c} \cos\left(2\pi f_c t + \frac{(2i-1)\pi}{4}\right) \quad 0 \le t \le T_s$$
 $i = 1, ..., 4$

$$s_i(t) = \frac{2E}{T_s} \cos\left(\frac{(2i-1)\pi}{4}\right) \cos(2\pi f_c t) - \frac{2E}{T_s} \sin\left(\frac{(2i-1)\pi}{4}\right) \sin(2\pi f_c t) \quad 0 \le t \le T_s$$

Signal space by two basis function

$$\phi_1(t) = \frac{2}{T_s} \cos(2\pi f_c t) \quad 0 \le t \le T_s$$

$$\phi_1(t) = \frac{2}{T_s} \sin(2\pi f_c t) \quad 0 \le t \le T_s$$



Probability of error in 4-QAM

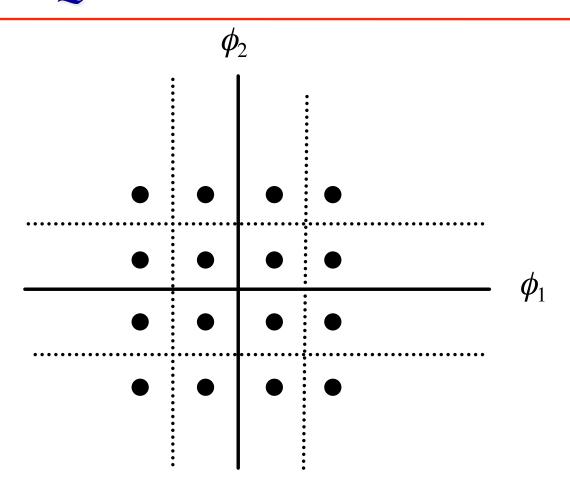
Minimum distance between two symbols $d_{\min} = \sqrt{2E}$

2 neighbors at minimum distance

$$P_e \simeq 2Q \left(\sqrt{\frac{E}{5N_0}}\right) = 2Q \left(\sqrt{\frac{2E_b}{N_0}}\right)$$



A more complex modulation: 16-QAM



More points can be added, this increases the transmit bit rate, but increases also the probability of error

Comparison of probabilities or error

$$P_{e,\text{BPSK}} = Q\left(\sqrt{\frac{2E}{N_0}}\right)$$

$$SNR = \frac{E}{N_0}$$

$$P_{e,4-\text{PAM}} \simeq 1.5Q \left(\sqrt{\frac{2E}{5N_0}}\right)$$

$$P_{e,4-\text{QAM}} \simeq 2Q \left(\sqrt{\frac{E}{N_0}}\right)$$

- Basic of modulation theory
- Probability of error over AWGN channels
 - BPSK
 - Amplitude Modulation
- Probability of error over fading + AWGN channels

Communication over wireless channel

 So far, we have seen only channels where the transmitted signal is received corrupted by AWGN, but in real channel

$$r(t) = \sqrt{A}s(t) + n_0(t)$$

- The power of the transmitted signals is attenuated by the wireless channel
- How the probability of error is affected by the channel?

Probability of error over fading channels

 Consider a AWGN + Rayleigh channel with path loss, Rayleigh fast fading, and fixed shadow fading

$$P_r = P_t G_t(\theta_t, \psi_t) G_r(\theta_r, \psi_r) \frac{\lambda^2}{(4\pi r)^2} \bar{PL} \cdot y \cdot z$$

$$\triangleq P_t C \cdot z$$



Probability of error over fading channels

The receiver sees the transmit power $P_t = \frac{E}{T_s}$ received with an attenuation $C \cdot z$ and thus

$$SNR = \frac{E}{N_0}Cz$$

 To compute the probability, mathematically, it is AS if the transmit power WERE

$$P_t = \frac{E}{T_s} Cz$$

and the channel WERE just AWGN and no fading

 Thus, the probability of error with fading has same expression of simple AWGN channels, but with the new SNR above

$$P_{e,\mathrm{BPSK}} = Q\left(\sqrt{\frac{2E}{N_0}}\right)$$
 \longrightarrow $P_{e,\mathrm{BPSK}}(z) = Q\left(\sqrt{\frac{2ECz}{N_0}}\right)$ AWGN+ Rayleigh fading

Probability of error over fading channels

- The probability so derived over fading channel is instantaneous
 - Depends on the given realization of the fading channel z
- What is the average probability of error, where the average is taken over the distribution of the fading?
- Just take the expectation of $P_{e,\mathrm{BPSK}}(z)$ over the distribution of z
- Thus (remember that the square of a Rayleigh random variable is an exponential random variable)

$$p(z) = \frac{1}{\gamma^*} e^{-\frac{z}{\gamma^*}}$$
 $\gamma^* \triangleq \mathbb{E} \, \text{SNR} = \frac{E}{N_0} C$ $\mathbb{E} \, z = 1$

$$\bar{P}_{e,\text{BPSK}} = \int_0^\infty P_{e,\text{BPSK}}(z)p(z)dz = \frac{1}{2} \left[1 - \sqrt{\frac{\gamma^*}{1+\gamma^*}} \right] \simeq \frac{1}{4\gamma^*}$$

Comparison of probability of error over fading channels

Simple AWGN channel:

$$P_{e,\text{BPSK}} = Q\left(\sqrt{\frac{2E}{N_0}}\right)$$

- linear increase of SNR results in exponential decrease in the error probability
- AWGN + Rayleigh channel:

$$\bar{P}_{e,\text{BPSK}} \simeq \frac{1}{4\gamma^*} \qquad \gamma^* \triangleq \mathbb{E} \,\text{SNR} = \frac{E}{N_0} C$$

- linear increase of the SNR gives only a linear decrease of error probability
- 40dB higher SNR than the simple AWGN channel to have same probability $P_e = 10^{-6}$

Packet error probability

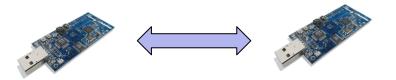
- Bits are grouped in unit denoted physical layer data frame, "physical layer message"
- What is the probability that the message is erroneously received?
- Suppose that the message is formed by L bits and BPSK is used, then the probability that the packet is in error is bounded by

$$P_m \le 1 - (1 - P_{e, BPSK})^L$$

Conclusions

Application
Presentation
Session
Transport
Routing
MAC

Phy



- We studied how bits are transmitted over the wireless channel by modulation and coding
- The probability of successful reception of messages was characterized for AWGN and for AWGN + fading channels

- But how a node has the right to transmit a message?
- Medium access control
 - When a node gets the right to transmit?
 - What is the mechanism to get such a right?