

## Lecture 5

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## Coprimality (Relative Primality)

**Definition.** Two integers  $m$  and  $n$  are coprime if their greatest common divisor is 1.

**Fact.** If  $a$  and  $n$  are coprime, then there exists a  $b$  such that  $ab = 1 \pmod n$ .

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**Excercise:** Why is this so?

## Chinese Remainder Theorem (CRT)

**Theorem.** (Sun Tzu 400 AC) Let  $n_1, \dots, n_k$  be positive pairwise coprime integers and let  $a_1, \dots, a_k$  be integers. Then the equation system

$$x = a_1 \pmod{n_1}$$

$$x = a_2 \pmod{n_2}$$

$$x = a_3 \pmod{n_3}$$

$$\vdots$$

$$x = a_k \pmod{n_k}$$

has a unique solution in  $\{0, \dots, \prod_i n_i - 1\}$ .

## Constructive Proof of CRT

1. Set  $N = n_1 n_2 \cdot \dots \cdot n_k$ .
2. Find  $r_i$  and  $s_i$  such that  $r_i n_i + s_i \frac{N}{n_i} = 1$  (Bezout).
3. Note that

$$s_i \frac{N}{n_i} = 1 - r_i n_i = \begin{cases} 1 & (\text{mod } n_i) \\ 0 & (\text{mod } n_j) \end{cases} \quad \text{if } j \neq i$$

4. The solution to the equation system becomes:

$$x = \sum_{i=1}^k \left( s_i \frac{N}{n_i} \right) \cdot a_i$$

# The Multiplicative Group

The set  $\mathbb{Z}_n^* = \{0 \leq a < n : \gcd(a, n) = 1\}$  forms a group, since:

- ▶ **Closure.** It is closed under multiplication modulo  $n$ .
- ▶ **Associativity.** For  $x, y, z \in \mathbb{Z}_n^*$ :

$$(xy)z = x(yz) \pmod n .$$

- ▶ **Identity.** For every  $x \in \mathbb{Z}_n^*$ :

$$1 \cdot x = x \cdot 1 = x .$$

- ▶ **Inverse.** For every  $a \in \mathbb{Z}_n^*$  exists  $b \in \mathbb{Z}_n^*$  such that:

$$ab = 1 \pmod n .$$

# Lagrange's Theorem

**Theorem.** If  $H$  is a subgroup of a finite group  $G$ , then  $|H|$  divides  $|G|$ .

**Proof.**

1. Define  $aH = \{ah : h \in H\}$ . This gives an equivalence relation  $x \approx y \Leftrightarrow x = yh \wedge h \in H$  on  $G$ .
2. The map  $\phi_{a,b} : aH \rightarrow bH$ , defined by  $\phi_{a,b}(x) = ba^{-1}x$  is a bijection, so  $|aH| = |bH|$  for  $a, b \in G$ .

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**Excercise:** How does this follow from CRT?

## Fermat's and Euler's Theorems

**Theorem.** (Fermat) If  $b \in \mathbb{Z}_p^*$  and  $p$  is prime, then  $b^{p-1} = 1 \pmod{p}$ .

**Theorem.** (Euler) If  $b \in \mathbb{Z}_n^*$ , then  $b^{\phi(n)} = 1 \pmod{n}$ .

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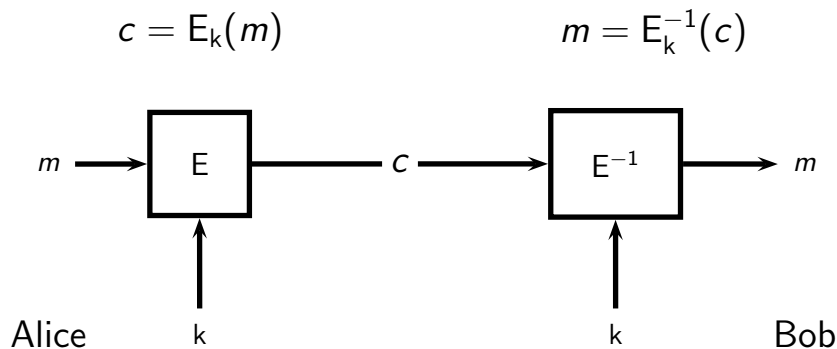
**Proof.** Note that  $|\mathbb{Z}_n^*| = \phi(n)$ .  $b$  generates a subgroup  $\langle b \rangle$  of  $\mathbb{Z}_n^*$ , so  $|\langle b \rangle|$  divides  $\phi(n)$  and  $b^{\phi(n)} = 1 \pmod n$ .

# Multiplicative Group of a Prime Order Field

**Definition.** A group  $G$  is called **cyclic** if there exists an element  $g$  such that each element in  $G$  is on the form  $g^x$  for some integer  $x$ .

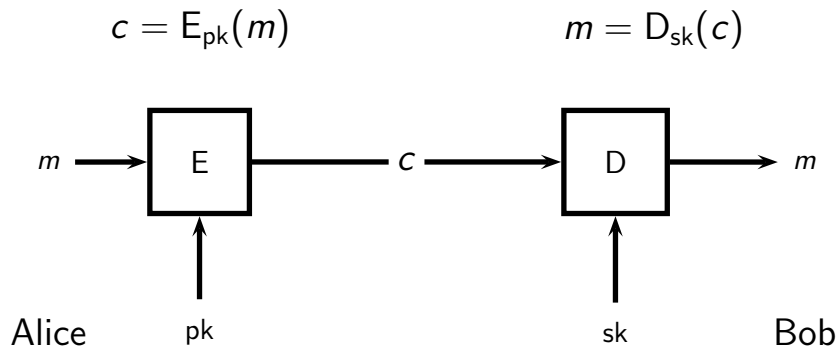
**Theorem.** If  $p$  is prime, then  $\mathbb{Z}_p^*$  is cyclic.

## Cipher (Symmetric Cryptosystem)





## Public-Key Cryptosystem



# History of Public-Key Cryptography

Public-key cryptography was discovered:

- ▶ By Ellis, Cocks, and Williamson at the Government Communications Headquarters (GCHQ) in the UK in the early 1970s (not public until 1997).
- ▶ Independently by Merkle in 1974 (Merkle's puzzles).
- ▶ Independently in its discrete-logarithm based form by Diffie and Hellman in 1977, and instantiated in 1978 (key-exchange).
- ▶ Independently in its factoring-based form by Rivest, Shamir and Adleman in 1977.

# Public-Key Cryptography

**Definition.** A public-key cryptosystem is a tuple  $(\text{Gen}, E, D)$  where,

- ▶  $\text{Gen}$  is a **probabilistic key generation algorithm** that outputs key pairs  $(pk, sk)$ ,
- ▶  $E$  is a (possibly probabilistic) **encryption algorithm** that given a public key  $pk$  and a message  $m$  in the plaintext space  $\mathcal{M}_{pk}$  outputs a ciphertext  $c$ , and
- ▶  $D$  is a **decryption algorithm** that given a secret key  $sk$  and a ciphertext  $c$  outputs a plaintext  $m$ ,

such that  $D_{sk}(E_{pk}(m)) = m$  for every  $(pk, sk)$  and  $m \in \mathcal{M}_{pk}$ .